KEMERBURGAZ UNIVERSITY INSTITUTE OF SCIENCE/ELECTRIC AND COMPUTER ENGINEERING DPARTMENT PHD PROGRAM



PROJECT NAME: Real Time CPU Scheduling

REQUIREMENT OF ECE 519 Advanced Operating Systems COURSE PREPEARED BY: AHMED M. ABDALLA MOHAMED STUDENT NO: 163106002

REAL TIME CPU SCHEDULING

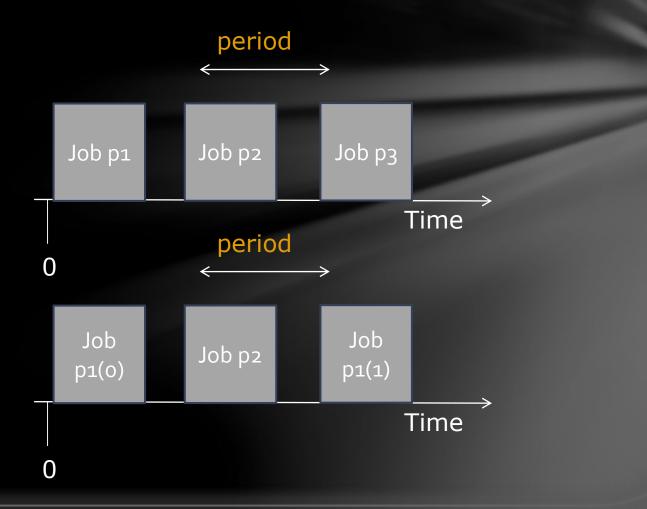
OUTLINES

- -> Basic Concepts.
 - Schedule. Real Time System. CPU burst VS I/O burst. CPU Scheduler. Dispatcher.
- -> Scheduling Criteria.
- Scheduling Algorithms.
 First-Come, First-Served (FCFS) Scheduling.
 Shortest-Job-First (SJF) Scheduling.
 Priority Scheduling.
 Round Robin (RR).

BASIC CONCEPTS

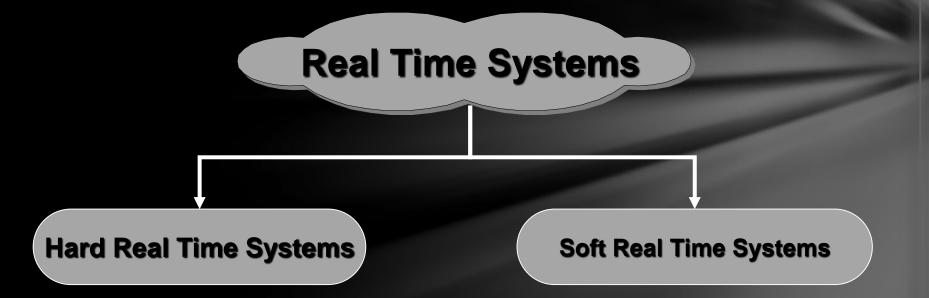
What is a Schedule?

Times at which process execute.



Real Time System:

Systems that performs critical tasks.



- Hard real time systems: critical tasks must be performed on time.
- Soft real time systems: critical tasks gets priority over other tasks.

CPU burst VS I/O burst

A CPU bursts when it is executing instructions

An I/O system bursts when it services requests to fetch information.

• • •	
load store add store read from file	CPU burst
wait for I/O	≻ I/O burst
store increment index write to file	CPU burst
wait for I/O	├O burst
load store add store read from file	<pre> CPU burst </pre>
wait for I/O	≻ I/O burst
•)

CPU Scheduler

- Short-term scheduler selects from among the processes in ready queue, and allocates the CPU to one of them

- Queue may be ordered in various ways
- CPU scheduling **decisions** may take place when a process:
 - 1. Switches from running to waiting state
 - 2. Switches from running to ready state
 - 3. Switches from waiting to ready
 - 4. Terminates
- Scheduling under 1 and 4 is **nonpreemptive**
- All other scheduling is preemptive

Dispatcher

- Dispatcher module gives control of the CPU to the process selected by the short-term scheduler, this involves:
- -> switching context
- > switching to user mode

- > jumping to the proper location in the user program to resume that program

 Dispatch latency – time it takes for the dispatcher to stop one process and start another running

SCHEDULING CRITERIA

- What do we want to achieve from scheduling (One of the following):

-> Max CPU utilization – keep the CPU as busy as possible

-> Max Throughput – of processes that complete their execution per time unit

-> Min Turnaround time – amount of time to execute a particular process

-> Min Waiting time – amount of time a process has been waiting in the ready queue

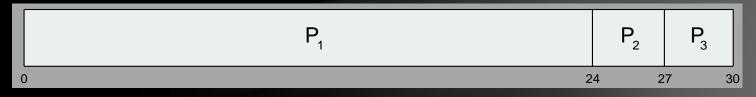
SCHEDULING ALGORITHMS

First-Come, First-Served (FCFS) Scheduling:

Example 1:- Suppose that the processes arrive in the order: P_1 , P_2 , P_3

<u>Process</u>	<u>Burst Time</u>
P ₁	24
P_{2}	3
P_3	3

The Gantt Chart for the schedule is:



• Waiting time for $P_1 = 0$; $P_2 = 24$; $P_3 = 27$

Average waiting time: (0 + 24 + 27)/3 = 17

(FCFS) Example 2:-

Suppose that the processes arrive in the order: P_2 , P_3 , P_1

The Gantt chart for the schedule is:



- Waiting time for $P_1 = 6$; $P_2 = o_1$; $P_3 = 3$
- Average waiting time: (6 + o + 3)/3 = 3
- Much better than previous case why ?

Shortest-Job-First (SJF) Scheduling: EXAMPLE:

Process	ArrivalTime	<u>Burst Time</u>
P ₁	0.0	6
P ₂	2.0	8
P ₃	4.0	7
P_4	5.0	3

• SJF scheduling chart

P ₄	P ₁	P ₃	P ₂
0 :	3) 1	6 24

Average waiting time = (3 + 16 + 9 + 0) / 4 = 7

Example of Shortest-remaining-time-first:

 Now we add the concepts of varying arrival times and preemption to the analysis.

Process	Arrival Time	<u>Burst Time</u>
P ₁	Ο	8
P_{2}	1	4
P ₃	2	9
P_4	3	5

• Preemptive SJF Gantt Chart

P ₁	P ₂	P ₄	P ₁	P ₃	
0	1 5	5 1	0 1	7	26

Average waiting time = [(10-1)+(1-1)+(17-2)+(5-3))]/4 = 26/4 = 6.5 msec

Priority Scheduling:

- A priority number (integer) is associated with each process
- The CPU is allocated to the process with the highest priority (smallest integer = highest priority)
 - Preemptive
 - Nonpreemptive
- Problem = Starvation low priority processes may never execute
- Solution = Aging as time progresses increase the priority of the process

Example of Priority Scheduling:

Process	Burst Time	<u>Priority</u>
P ₁	10	3
P ₂	1	1
<i>P</i> ₃	2	4
P ₄	1	5
P ₅	5	2

• Priority scheduling Gantt Chart

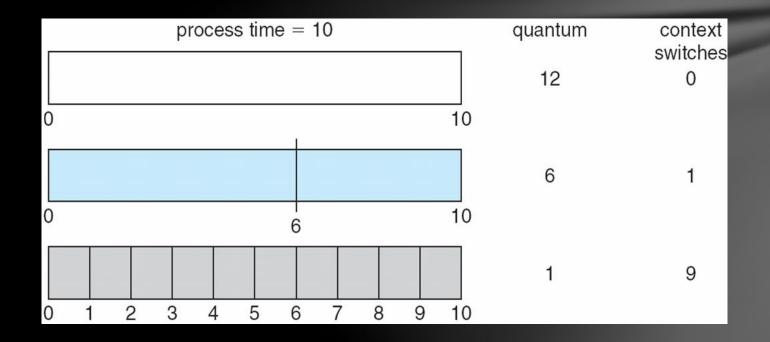
P ₂	P ₅	P ₁		P ₃	P ₄
0 1		6	16	18	3 19

• Average waiting time = 8.2 msec

Round Robin (RR):

- Each process gets a small unit of CPU time (time quantum q)
 - After this time has elapsed, the process is preempted and added to the end of the ready queue.
- If there are *n* processes in the ready queue and the time quantum is *q*,
 - each process gets 1/n of the CPU
 - No process waits more than (*n*-1)*q* time units.
- Timer interrupts every quantum to schedule next process Performance
 - If q is large \Rightarrow FIFO
 - $q \text{ small} \Rightarrow q \text{ must be large with respect to context switch time, otherwise overhead is too high}$
 - Context switch < 10 micro sec

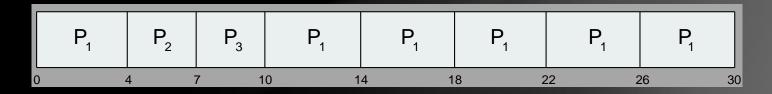
Time Quantum and Context Switch Time:



Example of RR with Time Quantum = 4

<u>Process</u>	<u>Burst Time</u>
P ₁	24
P ₂	3
<i>P</i> ₃	3

The Gantt chart is:



Better response than SJF

THANKYOU FORYOUR LISTENING

